Laboratory work # 1

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Problem # 1296 *Hyperjump*

Screenshot from Timus:



Explanation of algorithm:

Dynamic – *Kadane’s* algorithm:

Go through all elements, sum them, and save maximum value of sum. Negative sum is useless, which will be replaced with 0. The final output is the maximum of sum.

Computational complexity of algorithm:

Source code:

1. #include <iostream>
2. #include <cstdio>
4. **using** **namespace** std;
6. **int** main(){
8. **int** N;
9. **int** a[60001];
11. **int** ThisSum = 0;
12. **int** MaxSum = 0;
14. cin >> N;
15. **for**(**int** i = 0; i < N; i++){
16. cin >> a[i];
17. }
19. **for**(**int** i = 0; i < N; i++){
20. ThisSum += a[i];
21. **if**(ThisSum > MaxSum){
22. MaxSum = ThisSum;
23. }**else** **if**(ThisSum < 0){
24. ThisSum = 0;
25. }
26. }
27. cout << MaxSum << endl;
29. **return** 0;
30. }

Problem # 1155 *Troubleduons*

Screenshot from Timus:



Explanation of algorithm:

Categorize 8 cameras into two parts, namely *ACFH* and *BDEG*, the four cameras inside each section are not adjacent. If *A + C + F + H != B + D + E + G*, return impossible due to no solution. If the sum of two parts is equal, then we follow a principle, namely for adjacent destroys while for non-adjacent moving troubleduons between vertices of the same group, finally we can successfully move other cameras (*A, C, H*) in the same part to *F* (for instance), move other cameras (*D, E, G*) in the same part to *B* (For instance). Then the numbers of cameras *F* and *B* must be equal, so we can destroy two of them gradually. At last, we make it!

Computational complexity of algorithm:

Source code:

1. #include <cstdio>
2. #include <iostream>
3. **using** **namespace** std;
5. **int** a[8];
7. **void** move(**int** x,**int** y,**int** t)
8. {
9. **while**(a[x]){
10. **if**(a[t]==0){
11. a[t]++;a[y]++;
12. printf("%c%c+\n",t+'A',y+'A');
13. }
14. a[t]--;
15. a[x]--;
16. printf("%c%c-\n",t+'A',x+'A');
17. }
18. }
20. **int** main()
21. {
22. **int** i;
23. **for**(i=0;i<8;i++)
24. scanf("%d",&a[i]);
25. **int** s1=a[0]+a[2]+a[5]+a[7],s2=a[1]+a[3]+a[4]+a[6];
26. **if**(s1!=s2)
27. {
28. printf("IMPOSSIBLE\n");
29. **return** 0;
30. }
32. move(0,5,4);
33. move(2,5,6);
34. move(7,5,4);
35. move(3,1,0);
36. move(4,1,0);
37. move(6,1,2);
38. **while**(a[1]--)
39. printf("FB-\n");
41. **return** 0;
42. }

Problem # 1005 *Stone Pile*

Screenshot from Timus:



Explanation of algorithm:

The problem can be viewed as a 0-1 backpack problem by using dynamic programming to solve it. There are up to 20 stones, stones of varying weight, now is to take some stones as full as possible half of the total weight of the backpack, the rest is not full of another pile, so that the difference between the two piles can be minimized!

Computational complexity of algorithm:

Source code:

1. #include <iostream>
2. #include <cmath>
3. #include <algorithm>
4. **using** **namespace** std;
6. **int** main() {
7. **int** n;
8. cin >> n;
9. **int** w[n];
10. **for** (**int** i = 0; i < n; i++) {
11. cin >> w[i];
12. }
13. **int** sum = 0;
14. **for** (**int** i = 0; i < n; i++) {
15. sum += w[i];
16. }
17. **int** half = sum / 2;
18. **int** backpack[half+1];  //include  zero volume
19. // One-dimensional form given that two-dimentional form will exceed the space
20. **for** (**int** i = 0; i <= half; i++) {
21. backpack[i] = 0;
22. }
23. **for** (**int** i = 0; i < n; i++) {
24. **for** (**int** j = half; j >= w[i]; j--){
25. backpack[j] = max(backpack[j], backpack[j-w[i]] + w[i]);
26. }
27. }
28. /\* Prove:
29. A + B = Sum
30. A - B = Min
31. A = (Sum + Min) / 2
32. B = (Sum - Min) / 2
33. Min = abs(Sum - 2 \* min)
34. \*/
36. cout << abs(sum - 2 \* backpack[half]) << endl;
37. **return** 0;
38. }